

# Non-monotone Trust Region Technique for Equality Constrained Optimization

Xiao WU, Qinghua ZHOU\*

College of Mathematics and Information Science, Hebei University, Baoding, CHINA

**Abstract:** In this article, we propose and analyze a new trust region algorithm for solving equality constrained optimization problems. We incorporate a non-monotone strategy into trust region algorithm to construct a more relaxed trust region procedure and employ a differentiable exact penalty function. Under some reasonable conditions, the global convergence is established.

**Keywords:** Nonmonotone; Trust region methods; Constrained optimization; Exact penalty function; Global convergence

## 1. Introduction

In this article, we consider the following equality constrained optimization problem

$$\begin{aligned} \min_{x \in R^n} & f(x) \\ \text{s.t. } & h_i(x) = 0, i = 1, 2, \dots, m \end{aligned} \tag{1}$$

where  $f(x) : R^n \rightarrow R$  and  $h_i(x) : R^n \rightarrow R (i = 1, 2, \dots, m) (m \leq n)$  are assumed to be continuously differentiable functions.

Many authors have studied problem (1) (see [2,3,6,10-12]). These methods are monotonic algorithm. In 1982, Chamberlain in [1] proposed the watchdog technique for constrained optimization to overcome the Maratos effect. Inspired by this idea, Grippo, Lamparillo and Lucidi introduced a nonmonotone line search technique for Newton's method in [4]. Their conclusions were overall approachable for the nonmonotone method, especially when applied to highly nonlinear problems and in presence of narrow curved valley.

The nonmonotone methods are distinguished by the fact that they do not enforce strict monotonicity to the objective function values at successive iterations. Some researchers showed that utilizing non-monotone technique may improve both the possibility of finding the global optimum and the rate of convergence (see [13]). Due to the high efficiency of nonmonotone techniques, many authors are interested in working on employing nonmonotone strategies in various branches of optimization procedures (see [7,14]).

Although the nonmonotone technique has many advantages, it suffers from some drawbacks. Ahookhosh et al. introduced a modified nonmonotone strategy and employs it in a trust region framework in [8]. Their analysis of the new algorithm showed that it inherited both stability of trust region methods and effectiveness of the nonmonotone strategy.

In this paper we extend the nonmonotone technique [8] to trust region method for equality constrained optimization problems.

The rest of this paper is organized as follows: in Section 2, we describe a new nonmonotone trust region algorithm. In Section 3, we prove that the proposed algorithm is globally convergent. Finally, some conclusions are expressed in Section 4.

## 2. Algorithm

Before describing the new algorithm, we introduce some notations:  $g(x) = \nabla f(x), A(x) = \nabla h(x) (\nabla h_i(x), \nabla h_2(x), \dots, \nabla h_m(x)) \in R^{n \times m}$ . We define the matrix

$$\begin{aligned} P(x) &= I - A(x)(A(x)^T A(x))^{-1} A(x)^T \\ &= I - A(x)A(x)^+ \end{aligned} \tag{2}$$

where  $A(x)$  has full column rank and  $A(x)^+ = (A(x)^T A(x))^{-1} A(x)^T$ .

We know that a point  $x$  is called a stationary point of problem (1) if it satisfies the Kuhn-Tucker condition

$$P h(x) P^+ P P(x) g(x) P = 0 \tag{3}$$

Now we discuss our new nonmonotone trust region algorithm for solving problem (1). At  $k$ th iteration, if  $x_k$  does not satisfy the Kuhn-Tucker condition, we compute a trial step  $d_k$  by solving the following quadratic programming subproblem [2]

$$\begin{aligned} \min_{d \in R^n} & g_k^T d + \frac{1}{2} d^T B_k d \\ \text{s.t. } & h_k + A_k^T d \leq q_k \\ & P d \in \mathcal{P}_{\text{TR}}^* \end{aligned} \tag{4}$$

where  $B_k$  is an  $n \times n$  symmetric matrix which is the Hessian of the Lagrangian function at  $(x_k, l_k)$  or an

approximation to it,  $D_k > 0$  is a trust region radius,  $q_k$  is any number which satisfies

$$\min_{D_k \in [b_1, b_2]} Ph_k + A_k^T d_k P \# q_k \quad \min_{D_k \in [b_1, b_2]} Ph_k + A_k^T d_k P \quad (5)$$

and where  $b_1$  and  $b_2$  are two given constants that satisfy  $0 \# b_2 < b_1 < 1$ .

For testing whether the point  $x_k + d_k$  is accepted as the next iteration, we use the augmented Lagrangian merit function

$$F(x, l, s) = f(x) + l(x)^T h(x) + s Ph(x) P^2 \quad (6)$$

where  $l(x)$  satisfies

$$\min_{l \in R^m} Pg(x) - A(x)l P^2 \quad (7)$$

and  $s > 0$  is the penalty parameter.

Now, we define

$$\bar{R}_k = h_k F_{l(k)} + (1 - h_k) F_k \quad (8)$$

where

$$F_{l(k)} = \max_{0 \# j, m(k)} \{F(x_{k-j}, l_{k-j}, s_{k-j})\}$$

and  $0 \# m(k) \in \{m(k-1) + 1, N\}, m(0) = 0, N > 0, 0 \# h_{\min} < h_{\max} < 1$  and  $h_k \in [h_{\min}, h_{\max}]$ .

The actual reduction is

$$Ared_k = \bar{R}_k - F(x_k + d_k, l(x_k + d_k), s_k)$$

$$Pred_k = - (g_k + A_k l_k)^T d_k - \frac{1}{2} d_k^T B_k d_k - (l(x_k + d_k) - l_k)^T (h_k + A_k^T d_k) + s_k (Ph_k P^2 - Ph_k + A_k^T d_k P^2) \quad (9)$$

Therefore, the ratio is calculated

$$r_k = \frac{\bar{R}_k - F(x_k + d_k, l(x_k + d_k), s_k)}{Pred_k}$$

Now, we can outline our new nonmonotone trust region algorithm.

**Algorithm 1**

**Step 1** Given

$x_0 \in R^n, B_0 \in R^{n \times n}, D_0 > 0, e > 0, 0 \# m_1 < m_2 < 1, 0 < g_1 < g_2 < 1, 0 \# h_{\min} < h_{\max} < 1, N > 0, s_0 > 0, h > 0, 0 < b_2 < b_1 < 1$ . Set  $k = 0, m(0) = 0$ .

**Step 2** If  $Ph_k P + PP_k g_k P^2 \leq e$ , stop.

**Step 3** Solve the subproblem (4) to determine  $d_k$ . If  $d_k = 0$ , then stop; otherwise, calculate  $Pred_k$ . If

$$Pred_k \geq \frac{s_k}{2} (Ph_k P^2 - Ph_k + A_k^T d_k P^2) \quad (10)$$

does not hold, set

$$s_k = 2[(g_k + A_k l_k)^T d_k + \frac{1}{2} d_k^T B_k d_k + (l(x_k + d_k) - l_k)^T (h_k + A_k^T d_k)] / (Ph_k P^2 - Ph_k + A_k^T d_k P^2) + h \quad (11)$$

which ensures that the new value of expression (11) satisfies condition (10).

**Step 4** Compute  $Ared_k, Pred_k$  and  $r_k$ . If  $r_k \geq m_1$ , then set  $x_{k+1} = x_k + d_k$ .

**Step 5** Set

$$D_{k+1} = \begin{cases} [D_k, \bar{D}], & \text{if } r_k \geq m_2; \\ [g_2, D_k], & \text{if } m_1 \# r_k < m_2; \\ [g_1 D_k, g_2 D_k], & \text{if } r_k < m_1 \end{cases} \quad (12)$$

**Step 6** Update the matrix  $B_k$  to generate  $B_{k+1}$ .

Set  $s_{k+1} = s_k, k = k + 1$  and return to Step 2.

**3. Convergence Analysis**

In this paper, we consider the following assumptions that will be used to analyze the convergence properties of the new algorithm.

**Assumptions**

(H1) There exists a convex set  $W \subset R^n$  such that  $x_k, x_k + d_k \in W$  for all  $k$ .

(H2)  $f$  and  $h_i \in C^2(\cdot), i = 1, 2, L, m$ .

(H3) The matrix  $A(x) = \nabla h(x)$  has full column rank for all  $x \in W$ .

(H4)  $f(x), h(x), A(x), \nabla^2 f(x), (A(x)^T A(x))^{-1}$  and each  $\nabla^2 h_i(x), i = 1, 2, L, m$  are all uniformly bounded in norm in  $W$ .

(H5) The matrices  $\{B_k, k = 1, 2, L\}$  have a uniform upper bound, i.e. there exist  $b_1 > 0$  such that  $PB_k P \leq b_1$  for all  $k \in \{1, \dots, N\}$ .

In what follows, we introduce some basic Lemmas which play important role in the analysis of our new algorithm.

**Lemma 1.** Under the assumptions, there exists a positive constant  $b_2$  such that

$$Ph_k P^2 - Ph_k + A_k^T d_k P^2 \geq \min_{D_k \in [b_1, b_2]} h_k P, \frac{b_2 D_k}{PA_k^+ P} \quad (13)$$

$$Pred_k \geq \frac{1}{2} s_k Ph_k P \min_{D_k \in [b_1, b_2]} h_k P, \frac{b_2 D_k}{PA_k^+ P} \quad (14)$$

**Proof.** The inequality (13) can be found from Lemma 3.3 in [2]. The second result of the lemma is similar to Lemma 1 in [5].

**Lemma 2.** If Algorithm 1 does not terminate, then

$$(g_k + A_k l_k)^T d_k + \frac{1}{2} d_k^T B_k d_k \geq \frac{1}{4} PP_k \bar{g}_k P^2 \min_{D_k \in [b_1, b_2]} \frac{1}{PB_k P}, \frac{\bar{D}_k}{PP_k \bar{g}_k P} + PB_k PP_k d_k PP_k \bar{d}_k P$$

holds for all  $k$ , where

$$P_k = P(x_k) = I - A(x_k)(A(x_k)^T A(x_k))^{-1} A(x_k)^T \quad (15)$$

$$\bar{d}_k = (I - P_k)d_k \quad (16)$$

$$\bar{g}_k = g_k + B_k \bar{d}_k \quad (17)$$

$$\bar{D}_k = (D_k^2 - P_k \bar{d}_k P_k)^{\frac{1}{2}} \quad (18)$$

**Proof.** The proof is similar to Lemma 3 in [5], we omit it for convenience.

**Lemma 3.** Under Assumption, if Algorithm 1 does not terminate, there exists a positive constant  $m_1$  such that the inequality

$$\begin{aligned} & \text{Pr ed}_k - \frac{s_k}{2} (Ph_k P^2 - Ph_k + A_k^T d_k P^2) \\ & \geq \frac{1}{4} P P_k \bar{g}_k P^2 \min \left\{ \frac{1}{PB_k P}, \frac{\bar{D}_k}{P P_k \bar{g}_k P} \right. \\ & \left. - m_1 P d_k P P h_k P + \frac{s_k}{2} P h_k P \min \left\{ h_k P, \frac{b_2 D_k}{P A_k^+ P} \right\} \right. \end{aligned}$$

hold for all  $k$ .

**Proof.** Using Lemma 1, 2 and (9), we obtain

$$\begin{aligned} & \text{Pr ed}_k - \frac{s_k}{2} (Ph_k P^2 - Ph_k + A_k^T d_k P^2) \\ & \geq \frac{1}{4} P P_k \bar{g}_k P^2 \min \left\{ \frac{1}{PB_k P}, \frac{\bar{D}_k}{P P_k \bar{g}_k P} \right. \\ & \left. - PB_k P P d_k P \bar{d}_k P - (l(x_k + d_k) - l_k)^T \right. \\ & \left. (h_k + A_k^T d_k) + \frac{s_k}{2} P h_k P \min \left\{ h_k P, \frac{b_2 D_k}{P A_k^+ P} \right\} \right. \end{aligned} \quad (19)$$

According to Lemma 1, (15) and (16)

$$\begin{aligned} & P \bar{d}_k P P A_k^+ d_k P \\ & = P(A_k^+)^T (h_k + A_k^T d_k - h_k) P \\ & \leq 2 P A_k^+ P P h_k P \end{aligned} \quad (20)$$

By Assumptions and (7), there exists a positive constant  $m_2$  such that the inequality

$$P l(x_k + d_k) - l_k P \geq m_2 P d_k P \quad (21)$$

holds for all  $k$ . Hence, from Lemma1, (19), (20) and (21), we have

$$\begin{aligned} & \text{Pr ed}_k - \frac{s_k}{2} (Ph_k P^2 - Ph_k + A_k^T d_k P^2) \\ & \geq \frac{1}{4} P P_k \bar{g}_k P^2 \min \left\{ \frac{1}{PB_k P}, \frac{\bar{D}_k}{P P_k \bar{g}_k P} \right. \\ & \left. - 2 P A_k^+ P P B_k P P d_k P P h_k P - m_2 P d_k P P h_k P \right. \\ & \left. + \frac{s_k}{2} P h_k P \min \left\{ h_k P, \frac{b_2 D_k}{P A_k^+ P} \right\} \right. \end{aligned}$$

Therefore, under assumptions, there exists  $m_1 > 0$  such that the result of Lemma 3 is true.

**Lemma 4.** Under Assumptions, there exist positive constants  $m_3$  and  $m_4$ , such that, on the iterations that satisfy the condition

$$P h_k P \geq m_3 \quad (22)$$

we have

$$\text{Pr ed}_k - \frac{s_k}{2} (Ph_k P^2 - Ph_k + A_k^T d_k P^2) \geq m_4 \quad (23)$$

**Proof.** The proof is similar to Lemma 5 in [5].

**Lemma 5.** Under the assumptions, there exists a positive constant  $m_5$  such that

$$P F_k - F_{k+1} - P \text{red}_k P \geq m_5 P d_k P^2. \quad (24)$$

**Proof.** Using (6), (9) and Taylor expansion, we have

$$\begin{aligned} & P F_k - F_{k+1} - P \text{red}_k P \\ & = |f_k + g_k^T d_k + \frac{1}{2} d_k^T B_k d_k - f_{k+1} + l_{k+1}^T (h_k + A_k^T d_k) \\ & \quad - l_{k+1}^T h_{k+1} + s P h_k + A_k^T d_k P^2 - s P h_{k+1} P^2| \\ & \leq \frac{1}{2} |d_k^T (B_k - \nabla^2 f(x_k - x_1 d_k)) d_k| \\ & \quad + \frac{1}{2} |d_k^T (\nabla^2 h(x_k - x_2 d_k)) l_{k+1} d_k| \\ & \quad + s P d_k P^2 (P \nabla^2 h(x_k - x_2 d_k) P P h_k P \\ & \quad + P \nabla^2 h(x_k - x_2 d_k) P P A_k P P d_k P \\ & \quad + \frac{1}{4} P \nabla^2 h(x_k + x_2 d_k) P^2 P d_k P^2) \end{aligned}$$

where  $x_1, x_2 \in (0, 1)$ .

By Assumptions, there exists a positive constant  $m_5$  such that

$$P F_k - F_{k+1} - P \text{red}_k P \geq m_5 P d_k P^2.$$

**Lemma 6.** (See Lemma 5 in [9]) Under the assumptions, if  $P P_k g_k P + P h_k P > 0$ , then there exists a integer  $k_0$  and a positive constant  $\bar{s}$  such that for all  $k \geq k_0, s_k \geq \bar{s}$ .

**Lemma 7.** Under the assumptions and there exists an infinite set  $N$ , we have

$$\lim_{k \in N} F_{l(k)} = \lim_{k \in N} F_{k+1} = \lim_{k \in N} \bar{R}_k.$$

**Proof.** Using definition of  $\bar{R}_k$  and  $F_{l(k)}$ , we observe that

$$\begin{aligned} \bar{R}_k &= h_k F_{l(k)} + (1 - h_k) F_k \\ \bar{R}_{l(k)} &= h_{l(k)} F_{l(l(k))} + (1 - h_{l(k)}) F_{l(k)} \end{aligned}$$

And from the definition of  $F_{l(k)}$ , we have  $F_k = h_k F_{l(k)} + (1 - h_k) F_k$ ,

for any  $k \in N$ . Hence

$$\begin{aligned} F_k &= h_k F_{l(k)} + (1 - h_k) F_k \\ \bar{R}_{l(k)} &= h_{l(k)} F_{l(k)} + (1 - h_{l(k)}) F_k = \bar{R}_k. \end{aligned}$$

Then we have

$$F_k \neq \bar{R}_k \neq F_{l(k)}.$$

This fact, along with Lemma 4.7 in [7], leads us to have the conclusion.

**Theorem 1.** Under the assumptions. If Algorithm 1 fails to satisfy the termination condition, then

$$\lim_{k \rightarrow \infty} P h_k P = 0 \tag{25}$$

**Proof.** We prove it by contradiction. Assume that (25) is not true, that is, there exists a constant  $\bar{h} > 0$  such that

$$P h_k P \geq \bar{h}, \quad \forall k. \tag{26}$$

It follows from Lemma 1 and Step 4 that

$$\bar{R}_k - F_{k+1} \leq \frac{m_4 s_k}{2} P h_k P \min \left\{ \frac{b_2 D_k}{P A_k^+ P}, \frac{b_2 D_k}{P A_k^+ P} \right\} \tag{27}$$

If  $P h_k P \geq m_3$ , by Lemma 4 and (27), we have

$$\bar{R}_k - F_{k+1} \leq m_4 m_3 \tag{28}$$

Otherwise, (27) implies that

$$\bar{R}_k - F_{k+1} \leq \frac{m_4 s_k}{2} m_3^2 \tag{29}$$

Hence (28) and (29) imply

$$\bar{R}_k - F_{k+1} \leq m_4 \min \left\{ m_3, \frac{m_4 s_k}{2} m_3^2 \right\}, \quad \forall k \tag{30}$$

Using Lemma 7 and (30), we obtain

$$\lim_{k \rightarrow \infty} D_k = 0, \quad \lim_{k \rightarrow \infty} P d_k P = 0 \tag{31}$$

without loss of generality, it is assumed by

$$D_k \leq \frac{\bar{A}}{b_2} \bar{h}, \quad \forall k \tag{32}$$

By (26), (31), (32), Lemma 3 and Lemma 5, we can deduce  $D_{k+1} \leq \frac{\bar{A}}{b_2} \bar{h}$  for all  $k$  sufficiently large, which is in contradiction with (31). Therefore, (25) holds for all  $k$ .

**Theorem 2.** (See Theorem 3.9 in [2]) Under the assumptions, Algorithm 1 produces iterates  $\{x_k\}$ , which satisfy

$$\liminf_{k \rightarrow \infty} (P h_k P + P P_k g_k P) = 0.$$

### 4. Conclusions

In this paper, we propose a new non-monotone trust region algorithm for solving equality constrained optimization problems. After we analyzed the properties of the new algorithm, the global convergence theory is proved. We believe that there is considerable scope for modifying and adapting the basic ideas introduced in this paper. In the near future, we would like to combine the new algorithm with line search algorithm in order to sufficiently

use the information which the algorithm has already derived.

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